



# Robust Strategies for Stochastic Multi-Agent Systems

Extended Abstract

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## ABSTRACT

The precise probabilities of stochastic systems are often partially unknown and may face perturbations. Finding a strategy in this setting is difficult, as it requires dealing with uncertainty on the system transitions while interacting with other agents. In this paper, we introduce the robust model checking problem for Multi-Agent Systems, in which agents play strategies that ensure the satisfaction of a specification is satisfied, even though the system probabilities are uncertain. We consider specifications in a variant of Alternating-time Temporal Logic with bounded memory.

## KEYWORDS

Stochastic Multi-Agent Systems; Probabilistic Model Checking; Logics for Strategic Reasoning; Robustness

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## 1 INTRODUCTION

One of the first and most important developments in using formal methods for reasoning about strategies in Multi-Agent Systems (MAS) is the Alternating-time Temporal Logic (ATL) [1], which contains strategic modalities expressing cooperation and competition among agents to achieve a goal. Some aspects of MAS, such as the unpredictable behavior of agents and the occurrence of random phenomena, are uncertain. These aspects can be measured based on experiments or past observations and represented with stochastic models, such as Markov decision processes (MDP) and stochastic MAS. Probabilistic ATL (PATL) [8] extends ATL to the probabilistic setting, allowing reasoning about randomized strategic abilities of agents interacting in a system with stochastic transitions. Uncertainty in MAS may also originate from agents' partial observability of the system, but model-checking strategic abilities under

imperfect information and perfect recall entails undecidability, even when restricted to deterministic MAS [10] or to a single agent as in POMDPs [14], leading to consider memoryless or bounded-memory strategies [4, 5, 14].

In many cases, the precise probabilities of the system transitions are unknown and may face perturbations. An example is model-based reinforcement learning, where agents estimate the agent-environment interaction model (e.g., a MDP [13]). Since the model is learned from their interaction with the environment, its transitions are susceptible to error. Strategizing in such a setting requires dealing with uncertainty on the system transitions while interacting with other agents, who may be cooperative or adversarial. Different approaches exist: on Markov chains, MDPs, and POMDPs, it has been proposed to consider intervals for the possible value of transition probabilities [7, 11, 14]: strategies must hold for any system whose transition set is within the interval. A way to generalize uncertainty and robustness is by considering parameters, where transition probabilities can be represented as equations over a given set of parameters [2, 12].

We introduce the *robust model checking* problem, which ensures that a temporal specification is satisfied, even though the system probabilities may suffer perturbations. By relying on PATL-like specifications, we consider *coalitional* strategies, and capture the strategic behavior of agent coalitions in probabilistic MAS with an additional uncertainty concerning the exact transition probabilities.

## 2 PROBABILISTIC ATL WITH BOUNDED STRATEGIES

### 2.1 Parametric Systems

We propose to extend stochastic multi-agent systems into *parametric systems*. In a parametric system  $\mathcal{G}$ , transition probabilities are replaced with equations over a set of variables  $X$ . A set of parameters resulting in all probabilities being in  $[0, 1]$  is *well-defined*. Once a well-defined set of parameters  $Val$  is fixed, we obtain  $\mathcal{G}[Val]$ , a classical stochastic MAS. We thus consider if a property holds for all valuations over  $X$  that yield a well defined set of probabilities.

We follow Definition 10.97 of [3] of strategies: a *general randomized strategy*  $\sigma$  with bounded recall  $n$  is a tuple  $(Q, act, \Delta, start)$  where  $Q$  is a set of modes (or memory states),  $\Delta$  is a randomized transition function, *act* randomly selects the next action depending on the current state, and *start* randomly selects a starting mode for the strategy. As in [9], a strategy has finite memory  $n$  if  $|Q| = n$  and



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memoryless if  $|Q| = 1$ . A strategy is *observation-based* if  $\Delta$  and  $\text{act}$  can be represented as depending on  $2^{\text{AP}}$  and  $Q$ . This corresponds to *imperfect information*, since strategies may only depend on AP, the observed atomic propositions labeling the states, and not the states themselves. Let  $\text{Str}_{a,n}$  be the set of observation-based strategies with bounded recall  $n$  for agent  $a$  and  $\text{Str}_n = \cup_{a \in \text{Ag}} \text{Str}_{a,n}$ .

## 2.2 The Logic $\text{PATL}_b$

We introduce the Probabilistic Alternating-Time Temporal Logic with bounded strategies ( $\text{PATL}_b$ ) defined as follows:

**Definition 1.** The syntax of  $\text{PATL}_b$  is defined by the grammar

$$\varphi ::= p \mid \varphi \vee \varphi \mid \neg \varphi \mid \langle\langle C \rangle\rangle_k^{\bowtie d} (\mathbf{X}\varphi) \mid \langle\langle C \rangle\rangle_k^{\bowtie d} (\varphi \mathbf{U}\varphi) \mid \langle\langle C \rangle\rangle_k^{\bowtie d} (\varphi \mathbf{R}\varphi)$$

where  $p \in \text{AP}$ ,  $k \in \mathbb{N}$ ,  $C \subseteq \text{Ag}$ ,  $d$  is a rational constant in  $[0, 1]$ , and  $\bowtie \in \{\leq, <, >, \geq\}$ . Most of these operators are classical, except for  $\langle\langle C \rangle\rangle_k^{\bowtie d} \varphi$ , that asserts that there exists an observation-based strategy with complexity at most  $k$  for the coalition  $C$  to collaboratively enforce  $\varphi$  with a probability in relation  $\bowtie$  with constant  $d$ .

Given a coalition strategy  $\sigma_C \in \prod_{a \in C} \text{Str}_a^p$ , the set of possible outcomes of  $\sigma_C$  from a path  $\pi$  (i.e., a finite sequence of states) to be the set  $\text{out}_C(\sigma_C, \pi) = \{\text{out}((\sigma_C, \sigma_{-C}), \pi) : \sigma_{-C} \in \prod_{a \in \text{Ag}_{-C}} \text{Str}_a\}$  of probability measures that the players in  $C$  enforce when they follow the strategy  $\sigma_C$ , namely, for each  $a \in \text{Ag}$ , player  $a$  follows strategy  $\sigma_a$  in  $\sigma_C$ . We use  $\mu_{\pi}^{\sigma_C}$  to range over the measures in  $\text{out}_C(\sigma_C, \pi)$ .

**Definition 2.**  $\text{PATL}_b$  formulas are interpreted in a stochastic system  $\mathcal{G}$  and a path  $\pi$ . Most of the semantics is classical, except for the coalition operator, defined as follows<sup>1</sup>:

$$\mathcal{G}, \pi \models \langle\langle C \rangle\rangle_k^{\bowtie d} \varphi \quad \text{iff } \exists \sigma_C \in \prod_{a \in C} \text{Str}_{a,k} \text{ s.t. } \forall \mu_{\pi_0}^{\sigma_C} \in \text{out}_C(\sigma_C, \pi_0),$$

$$\mu_{\pi_0}^{\sigma_C}(\{\pi' : \mathcal{G}, \pi' \models \varphi\}) \bowtie d$$

## 3 THE ROBUST $\text{PATL}_b$ MODEL CHECKING PROBLEM

We introduce the model checking problem for  $\text{PATL}_b$ . Three cases are of interest: (i) In the most general case, we have an arbitrary *parametric system*. (ii) When every transition may be perturbed at most a fixed  $\epsilon$ , we have *interval perturbations* on the system. (iii) When only a few critical components have an uncertain behavior, we can assume the number of perturbed transitions is *fixed*.

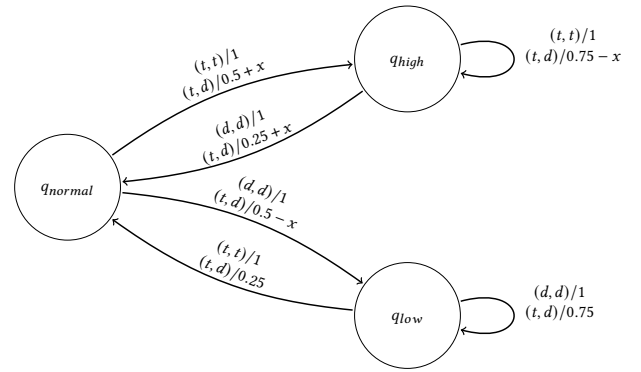
**Definition 3.** For a parametric system  $\mathcal{G}$ , state  $s \in \text{St}$ , and formula  $\varphi$  in  $\text{PATL}_b$ , the parametric model checking problem for  $\text{PATL}_b$  consists in deciding if for all well-defined valuations,  $\mathcal{G}[\text{Val}]$ ,  $s \models \varphi$ .

**Example 1.** Let us consider a system  $\mathcal{G}_{\text{river}}$  with two companies sharing the usage of a river. At every step, each company has two available actions: discharge wastewater directly into the river (action  $d$ ) or treat it before discharging it into the river (action  $t$ ). Atomic propositions state whether the river's water quality has reached low (proposition  $\text{low}$ ) or high levels (proposition  $\text{high}$ ). The system is shown in Figure 1. The propositions  $\text{low}$  and  $\text{high}$  are true only in state  $q_{\text{low}}$  and  $q_{\text{high}}$ , resp. In state  $q_{\text{normal}}$ , no proposition is true, representing that the water quality is normal. If both companies discharge the

<sup>1</sup>See [6] for the complete definition.

wastewater, the water quality is guaranteed to decrease (from high to normal, and normal to low). Similarly, if both companies treat the water, the quality will increase (from low to normal, and normal to high). When only one company treats the water while the other discharges the wastewater, the effect has a degree of uncertainty.

When in  $q_{\text{normal}}$ , the probability that the water quality increases when only one company treats the water may be slightly higher than planned, and so the probability to go to  $q_{\text{high}}$  is  $0.5 + x$ . As a consequence, the probability to go to  $q_{\text{low}}$  under this situation is  $0.5 - x$ , to compensate. At the same time, the probability to stay in  $q_{\text{high}}$  with only one company treating water may be slightly lower than expected, but follows the same trend, hence it is  $0.75 - x$ . Still, we can check that, as long as  $x$  stays within 0 and 0.25, a company alone can always make sure to have at least probability 0.375 of having a good water quality within two time steps by treating the water.



**Figure 1:** The parametric system  $\mathcal{G}_{\text{river}}^p$  where some transitions may change together depending on  $x \in [0, 0.25]$ .

## 4 CONCLUSION

This paper introduces the intricate problem of verifying the robustness of strategies for agents operating within stochastic multi-agent systems (MAS) that are prone to perturbations or variations. We introduced  $\text{PATL}_b$ , a logic tailored for reasoning about observation-based bounded memory strategies, which we model as automata. In our context, perturbations are parameterized, and we explored two distinct cases: one where the number of parameters is fixed, and another where the perturbations can assume any value within a bounded interval.

The robust model-checking problem guarantees that strategies are resilient to various types of perturbations in models, which is crucial in applications where exact probabilities of random events are imprecise. Considering bounded memory allows agents to retain relevant information while avoiding the undecidability issues associated with the combination of perfect recall and imperfect information in ATL-based formalisms [10].

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